



Not One of Them

Costly Binary Search

Log Drivin' Hirin'

City Hall

Harbingers

Camel and Oases

Boxes

Shortsighted

Jumping Stones

Sending Blessings

Go To Goal

Healthy Lifestyle

Frequent Alphabet

Topic-related  
tasks

Still DP

“Stolen” from a  
contest, sorted by  
(expected)  
hardest to easiest

I	Not One of Them	Topic-related tasks
C	Costly Binary Search	
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F	Camel and Oases	Still DP
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A	Go To Goal	"Stolen" from a contest, sorted by (expected) hardest to easiest
H	Healthy Lifestyle	
G	Frequent Alphabet	

as promised, one more task  
that has not been discussed  
yesterday

IOI 2014 Holiday (reduced)

given  $N$  cities in a line. City  $i$  has value  $A[i]$ . You start at city 0. In one day, you can either move to neighbouring city or take the value (at most once) of the current city.

For each  $d=0..2N$ , determine the maximum total value you can get if you have  $d$  days

$$1 \leq N \leq 100k$$



find the value of single  $d$  can be done in  $O(N \lg N)$  by iterating which rightmost city to be visited.

let  $\text{opt}(d)$  = the farthest city you visit when you have  $d$  days in the optimal solution.

we have  $\text{opt}(d) \leq \text{opt}(d+1)$

# Graph Connectivity

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let's practice some  
tasks

# IOI 2015 Practice Graph

given a graph and two nodes A, B.  
determine how many vertices which, if  
removed, disconnects A and B

$$1 \leq N \leq 100.000$$

$$0 \leq M \leq 200.000$$



run DFS tree with root = A

the possible candidate vertices  
are vertices in the path from A to B

node  $u$  is an answer if

let's say  $v$  is a child of  $u$  where  
 $v = \text{ancestor}(B)$ .

$$\text{dfs\_low}[v] \geq \text{dfs\_num}[u]$$

more

[https://www.hackerearth.com/practice/  
algorithms/graphs/strongly-connected-  
components/practice-problems/algorithm/  
a-walk-to-remember-qualifier2/](https://www.hackerearth.com/practice/algorithms/graphs/strongly-connected-components/practice-problems/algorithm/a-walk-to-remember-qualifier2/)

given a directed graph, for each node, determine whether there is a cycle starting from the node

$$1 \leq N \leq 100.000$$

$$1 \leq M \leq 200.000$$



just check for each node  
whether that node is alone in  
the SCC

one more

[http://acm.timus.ru/forum/thread.aspx?  
id=22089&upd=633721365703625916](http://acm.timus.ru/forum/thread.aspx?id=22089&upd=633721365703625916)

given a directed graph, determine which nodes can go to ALL other nodes.

$$1 \leq N \leq 100.000$$

$$0 \leq M \leq 200.000$$



so the observation is, if  $a$  and  $b$  is  
in one SCC, then the set of  
vertices that can be visited by  $a$   
and  $b$  is exactly the same.

therefore, run SCC, group nodes in one SCC to be one node.

There is an edge from SCC node  $a$  to  $b$   $\Leftrightarrow$  there is an edge from node  $u$  to  $v$  where  $u$  is in  $a$  and  $v$  is in  $b$

this technique is quite common. let's name it SCC graph

now we got a DAG (otherwise SCC is not optimal)

then, just choose a candidate node (vertex without an indegree), then check whether that node can visit all other nodes

last

# ICPC Jakarta Regional 2012

## Unique Path

given a graph, find the number of pair  
of nodes with unique path

$$2 \leq N \leq 10k$$

$$1 \leq M \leq 100k$$

EOF

Q&A?